Mina: Economics and Monetary Policy

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Abstract

As cryptocurrencies become more popular and widely used, the cost to verify the blockchain grows proportionally with total transaction throughput and quickly becomes out of reach for most real-world users. In this way, cryptocurrencies can end up victims of their own success, becoming increasingly inaccessible to their users who in turn are required to trust intermediaries. Mina Protocol solves this problem by replacing the blockchain with an easily verifiable proof. Such a protocol requires a careful separation of roles, incentive design, and sound monetary policy to function properly and resist exploitation. We present the main considerations below.

1 Introduction

By combining cryptographic techniques with economic incentives, Bitcoin[Nak09] became the first decentralized peer-to-peer payments protocol. A critical aspect of Bitcoin's design is the extent to which different components reinforced and strengthened one another. For example, Proof of Work imposes a cost on block production, which discourages bad actors from submitting invalid transactions. However, that incentive is dramatically multiplied by the use of a blockchain to store historical transactions, because it makes any tampering immediately evident to all peers on the network with a very low computation cost. This technique—making attacks costly and defense cheap—is a cornerstone of modern cryptography and a hallmark of a well-designed cryptoeconomic protocol.

However, Bitcoin's design has some limitations. While blockchains enable rapid tamper-detection, they do not provide rapid evidence of correctness. In fact, each time a new participant joins the network, they must check every transaction since the beginning of the network to verify correctness. This requirement grows linearly with total transaction throughput and quickly becomes out of reach for most real-world users on resource constrained devices like smartphones. **In this way, as cryptocurrencies become more popular and widely used, they end up victims of their own success, becoming increasingly inaccessible to their users who in turn are required to trust intermediaries.**

Recursive zk-SNARKs provide a solution to this dilemma. Similar to how blockchains enable constant-time tamper-evidence of a historical ledger, recursive zk-SNARKs enable constant-time correctness-verification of a historical ledger. Instead of each participant on the network verifying historical transactions for themselves, the network collaborates to generate proofs-of-correctness for transactions (zk-SNARKs), and then shares those around the network. So instead of end users trusting intermediaries to provide accurate information about the state of the ledger, they are given the state along with a zk-SNARK which cryptographically guarantees that state is accurate. **In short, Mina Protocol replaces the blockchain with an easily verifiable proof.** For more details on the technical implementation of this protocol, refer to the technical whitepaper[BMRS20].

Such a protocol requires a careful separation of roles, incentive design, and sound monetary policy to function properly and resist exploitation. We describe the main design considerations below.

2 Network Roles and Incentives

Most cryptocurrency protocols have at least two roles in their networks: 1) those who verify every transaction in the network, often called full nodes, stakers, or miners, and 2) those who trust third parties to verify transactions for them, such as light clients. As these protocols gain adoption, it becomes increasingly expensive to verify the ledger, so more participants become priced out of the first group and pushed into the second. For example, even though Bitcoin has averaged less than 1.5 transactions per second historically, a new member of the network must verify nearly 500,000,000 transactions in order to have full-node security. This problem is compounded in some proposed cryptocurrencies, which claim to have 10x-100,000x the transaction throughput of Bitcoin, and therefore generate gigabytes or terabytes of data every week under peak loads.

In contrast, Mina has constant resource requirements: regardless of how many transactions the network has processed, users can fully verify the current state with just a small zk-SNARK. To support this, Mina has three roles in the network, each incentivized to participate by different mechanisms.

2.1 Verifiers

We expect the vast majority of network participants to be capable of verifying. Because Mina uses recursive zk-SNARKs to continuously prove state validity, full-node security is achieved by simply downloading a zk-SNARK, which is around a few hundred bytes and takes a few milliseconds of compute to verify. The zk-SNARK certifies consensus information and a Merkle root to a recent ledger state. At this point, verifiers can request Merkle paths to relevant parts of the state. By checking the Merkle path, verifiers ensure that the parts of the state they care about (such as their account balances) are indeed contained within the same ledger that's certified by the zk-SNARK.

2.2 Block Producers

Block producers are akin to miners or stakers in other protocols. They are incentivized by protocol distributions in the form of block rewards or coinbase transactions, as well as network fees paid by users. Importantly, block producers are not incentivized by the threat of slashing in order to participate, as Mina uses Ouroboros[DGKR17]. In addition to directly staking, individuals can delegate their stake to another block producer. This allows the delegate to stake—but not send transactions—on another's behalf.

As is common, block producers select which transactions to include in the next block. Obviously, they are incentivized to include the highest fee transactions. However, in order to ensure the blockchain remains succinct, block producers have an additional responsibility: for each transaction they add to a block, they must SNARK an equivalent number of previously added transactions. If they don't do so, their block won't comply with the consensus rules and will be rejected by other nodes. It may be helpful to imagine a queue of transactions. If a block producer wants to add 10 transactions to the back of the queue (so they can claim the transaction fees), they must SNARK 10 transactions from the front of the queue. They may produce those SNARKs themselves, or they may select them from a marketplace that other specialized network participants, Snarkers, contribute to.

2.3 Snarkers

Snarkers, described in the technical Mina whitepaper[BMRS20], are network participants who produce zk-SNARKs that verify transactions. They are compensated by posting fees for doing so, called bids, and if their SNARKs are used in a block, the block producer pays out those fees from the total transaction fees.

Given that many different Snarkers can post fees for a single transaction, and block producers are incentivized to minimize the fees they pay out for SNARK work, this naturally forms a marketplace where participants compete to produce the most cost-efficient zk-SNARK proofs. For convenience, we can refer to this as the Snarketplace. Below, we'll evaluate some of its economic dynamics.

First, as long as there is some amount of latent compute available on the network——almost certainly assured by a network of block producers——the Snarketplace will not affect the protocol's liveness or censorship resistance, since there will always be a bid to produce SNARKs at some cost.

To see this, consider the incentives of a network participant who has spare compute capacity and observes SNARK fees rising on the network due to a lack of SNARK production. When SNARK fees go above the marginal cost of SNARK production for this participant, if they are a rational economic actor, they will begin offering bids above their price of production and below market price to capture profit, therefore stabilizing SNARK fees. Block producers can simply pass along those fees to those who are transacting. In practice, the computation cost to generate a SNARK should be extremely low, in the range of less than a penny per transaction at current cloud computing prices.

Second, given economies of scale, it's possible that some SNARK operations could become more dominant over time, as some mining operations have with Bitcoin. However, any resulting concentration of SNARK production would not have any bearing on the censorship resistance or liveness of the protocol. This is because anyone will be able to produce SNARKs at some price. So even if there was some SNARK producer who could offer cheaper SNARKs, if they refused to process certain transactions, anyone else on the network could provide a SNARK for a slightly higher fee, and block producers would still be incentivized to include it in the ledger. Third, an explicit, voluntary coordination mechanism that prevents double-proving would likely reduce SNARK fees and computational waste on the network, though such a mechanism isn't necessary for security or liveness.
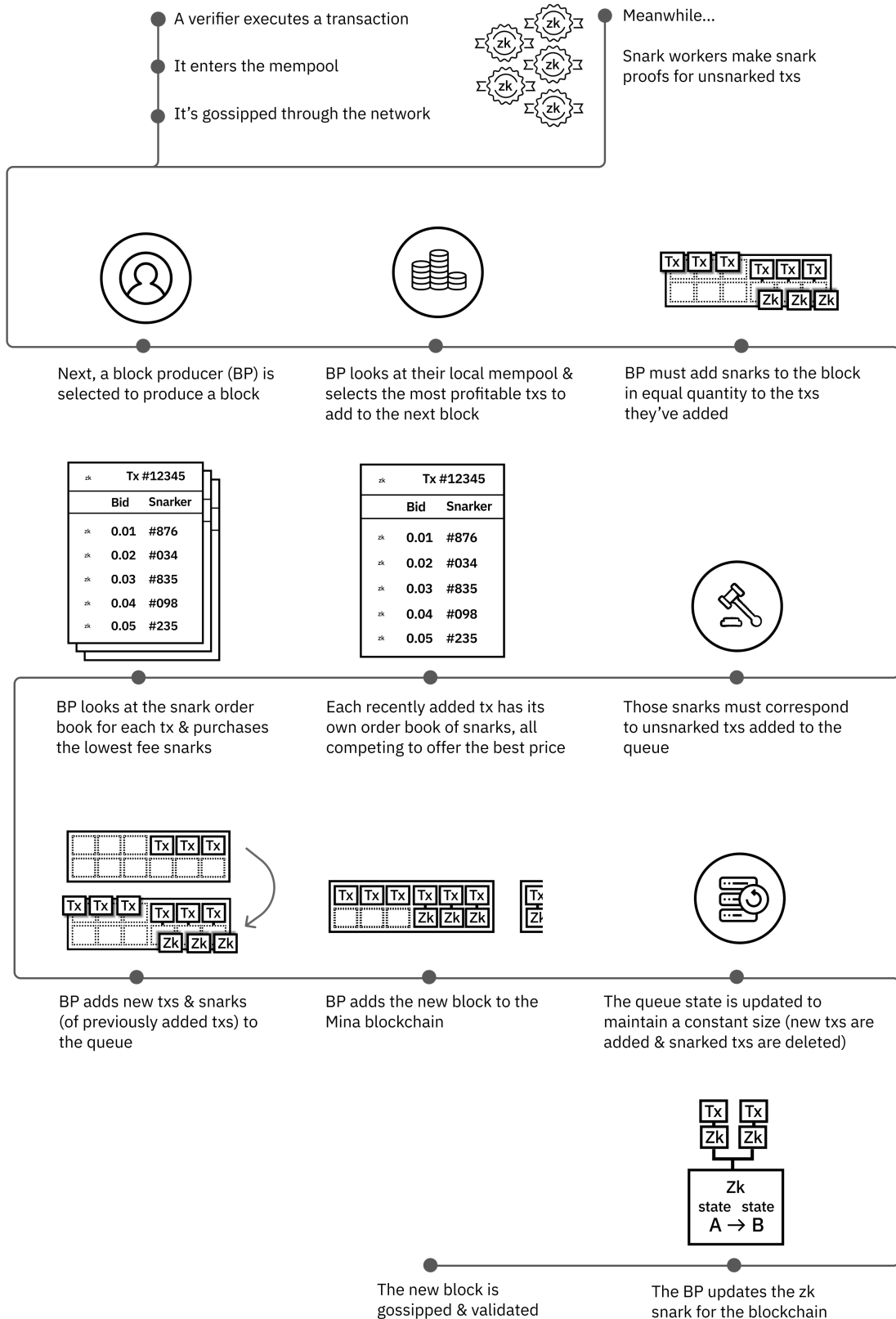
2.3.1 Censorship resistance

It may be worthwhile to expand a bit on the argument supporting censorship resistance. While there are a number of plausible attacks to bitcoin's censorship resistance[NBF+16], it turns out you can be included in a block if you just pay a high enough fee or wait long enough. Mina is very similar to bitcoin, with the added complication that you also need your transaction to be Snarked at some point after it is included in the block.

So for Mina to not be censorship resistant, a transaction would have to be included in a block and then never Snarked. That would require collusion among all Snarkers (who would refuse to Snark the transaction) or all block producers (who choose not to include the Snarked transaction in the ledger updates). By Ouroboros assumptions we can presume the network is decentralized and so block producers aren't colluding.

Now let's turn to Snarkers. Because the computation cost of Snarking a transaction is both extremely low and accessible to anyone with a computer, getting all Snarkers to collude is far more difficult than getting all block producers to collude. Existing Snarkers cannot prevent a new Snarker from entering the market. If someone notices a transaction not being Snarked, they could easily and profitably offer to Snark it, and even with very low transaction fees a block producer could profitably include the Snarked transaction in a ledger update.

Lifecycle of a Mina Transaction

A verifier executes a transaction

It enters the mempool

It's gossiped through the network

Meanwhile...

Snark workers make snark proofs for unsnarked txs

Next, a block producer (BP) is selected to produce a block

BP looks at their local mempool & selects the most profitable txs to add to the next block

BP must add snarks to the block in equal quantity to the txs they've added

zk	Tx #12345	
	Bid	**Snarker**
zk	**0.01**	**#876**
zk	**0.02**	**#034**
zk	**0.03**	**#835**
zk	**0.04**	**#098**
zk	**0.05**	**#235**

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BP looks at the snark order book for each tx & purchases the lowest fee snarks

Each recently added tx has its own order book of snarks, all competing to offer the best price

Those snarks must correspond to unsnarked txs added to the queue

BP adds new txs & snarks (of previously added txs) to the queue

BP adds the new block to the Mina blockchain

The queue state is updated to maintain a constant size (new txs are added & snarked txs are deleted)

Zk
state   state
A → B

The new block is gossipped & validated

The BP updates the zk snark for the blockchain

Economics of a Mina Transaction

Block Producer

1. Stakes Mina, is selected to produce a block
2. Selects the least expensive snark bids for unsnarked txs and adds them to the queue
3. Selects the least expensive snark bid & adds it to the queue
4. Earns the transaction fees
5. Pays the snarkers
6. Earns the block rewards and the spread between the snarker bids & tx fees

Anatomy of the Next Block
Block Production Queue

| | Snarked transaction | Unsnarked transaction | Unused transaction queue space |
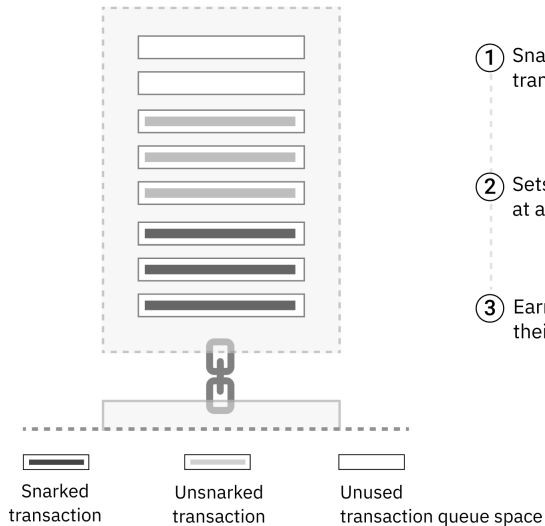
Snarker

1. Snarks recently added transactions
2. Sets bids to sell the snarks at a competitve rate
3. Earns what they bid if their snark is selected

Transaction Pool
Orderbook

1 Tx: 5 mina	✓
1 Tx: 4 mina	✓
1 Tx: 3 mina	✓
1 Tx: 3 mina	
1 Tx: 3 mina	
1 Tx: 2 mina	
1 Tx: 2 mina	
1 Tx: 2 mina	
1 Tx: 1 mina	

One Snark Per Transaction
Other block producers won't build on blocks that lack snark-to-transaction parity.

Multiple Bids Per Snark
The snarketplace offers a range of bids to snark each transaction.

Let the Market Decide
Block producers have natural incentive to choose the most profitable transactions and least expensive snarks.

Snarketplace
Orderbook
(per transaction)

1 Snark: 0.5 mina	
1 Snark: 0.4 mina	
1 Snark: 0.3 mina	
1 Snark: 0.3 mina	
1 Snark: 0.3 mina	
1 Snark: 0.2 mina	
1 Snark: 0.2 mina	
1 Snark: 0.2 mina	
✓ **1 Snark:** 0.1 mina	

# 3   Nomenclature and Denomination

Individual units of the currency are called mina. The initial supply of Mina is 1,000,000,000. Each mina is divisible into 1 billion units (nanomina). Intermediary denominations of the currency use standard metric prefixes:

1. 1 = 1 mina

2. 0.01 = 1 centimina (cent)

3. 10^{-6} = 1 micromina

4   Monetary Policy

Because Mina uses a variant of Ouroboros Proof of Stake[DGKR17][BMRS20], block rewards and fees are distributed approximately pro-rata to current holdings as long as everyone in the network is staking. Therefore, assuming high staking participation, any nominal inflation in the protocol is cancelled out by a nominal return denominated in the protocol, ensuring that accounts' proportional holdings of the currency remain constant.

However, those who do not choose to stake or delegate will experience dilution relative to those who do. In order to motivate participants to begin staking, Mina's nominal inflation will begin at 12%. Then, over the course of the first five years, the inflation rate will fall to 7%, and remain at 7% thereafter by default, subject to changes via the chain's governance[1].
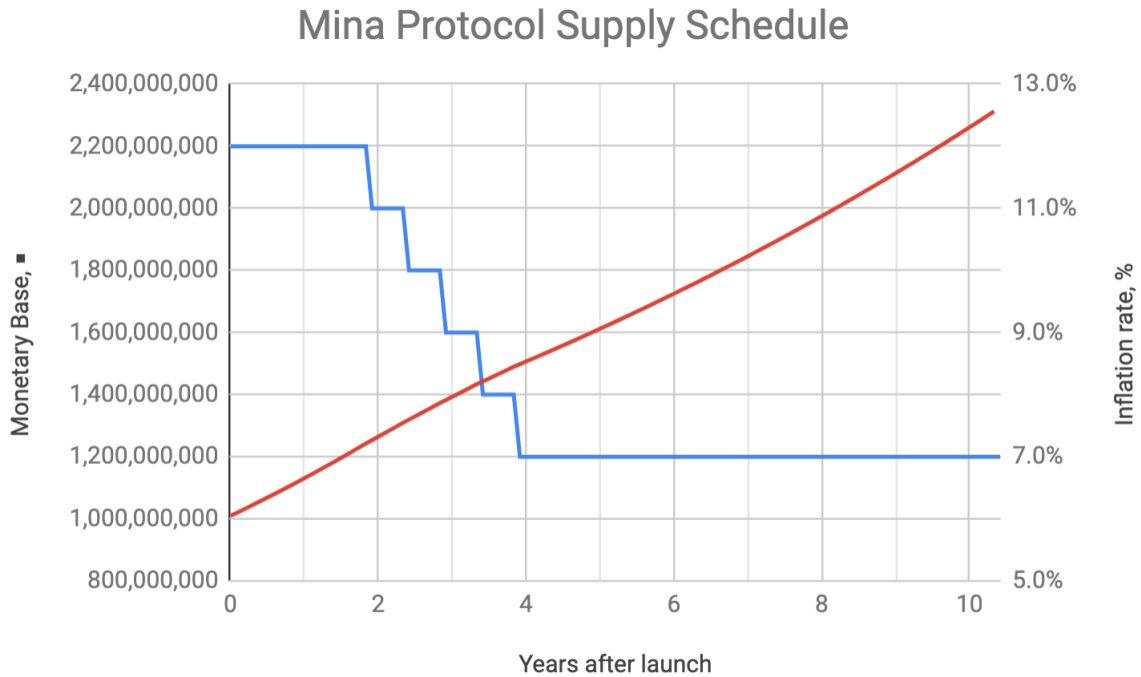
Importantly, the protocol will target these inflation rates regardless of the staking participation rate. That means that block rewards will change dynamically to target this inflation rate. For example, if only 50% of the network stakes, then block rewards will double. This is because per Ouroboros, number of blocks produced per epoch is expected to be proportional to the staking ratio. This naturally encourages more individuals to stake at low participation rates.

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[1]A governance paper is forthcoming.

Block Reward Multiplier by Participation Rate
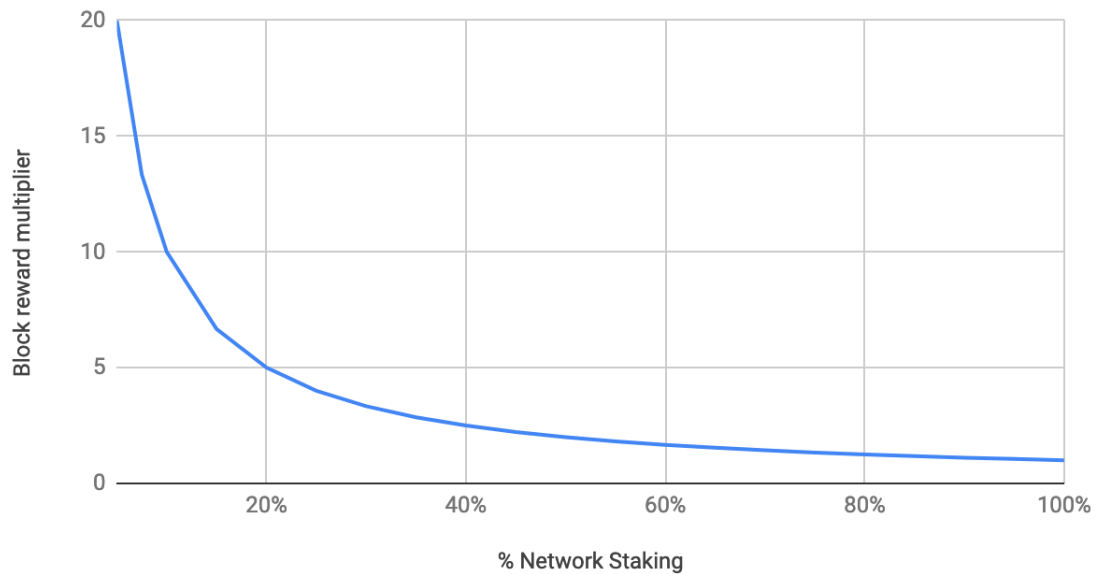
The primary consideration when lowering inflation rates would be determining whether rewards are higher than is necessary to ensure block producers have sufficient incentive to validate without requiring a large increase in fees.

It is also important to architect how monetary policy enables the ongoing development of a decentralized network such as Mina. Inflation, in the form of block rewards, is used to secure the chain against attacks by incentivizing validation by block producers. However, security is not the only requirement for a new network. Namely, long-term maintenance and improvement of the protocol will also require funding. We suggest the most suitable mechanism for this is the creation of special block rewards, allocated to recipients as determined by the chain's governance.

References

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[DGKR17] Bernardo David, Peter Gaži, Aggelos Kiayias, and Alexander Russell. Ouroboros praos: An adaptively-secure, semi-synchronous proof-of-stake protocol. Cryptology ePrint Archive, Report 2017/573, 2017. http://eprint.iacr.org/2017/573.

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